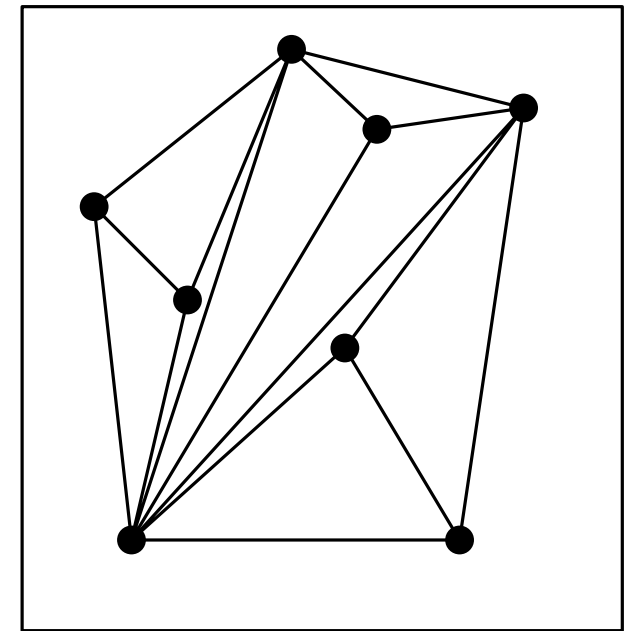
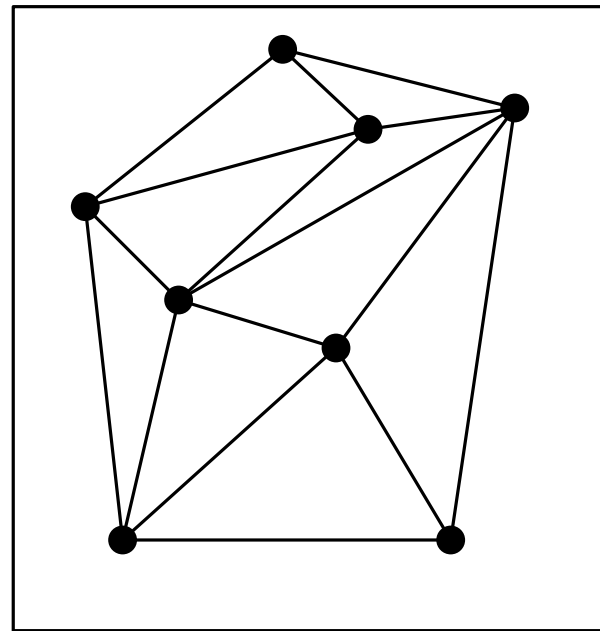
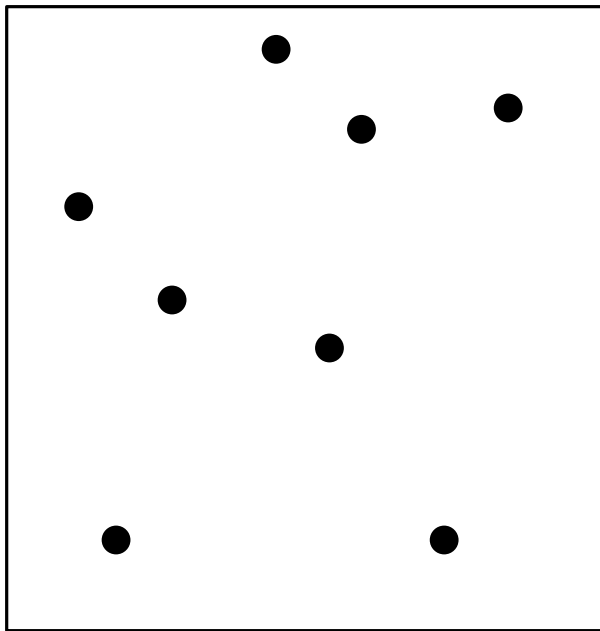


Counting and Enumeration in Combinatorial Geometry

Günter Rote
Freie Universität Berlin



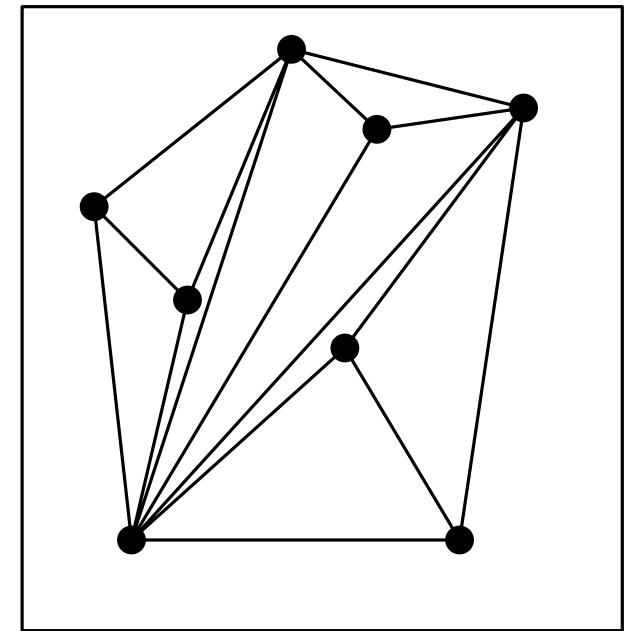
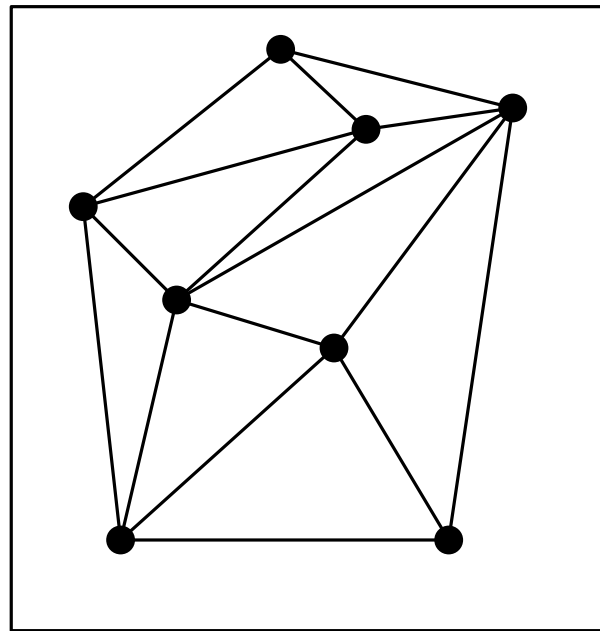
two triangulations

General position: No three points on a line

Counting and Enumeration in Combinatorial Geometry

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Freie Universität Berlin

- enumeration
- counting and sampling
- bounds
- optimization
- ...



two triangulations

General position: No three points on a line

Given a set of n points in the plane in general position,
how many

- triangulations
- non-crossing spanning trees
- non-crossing Hamiltonian cycles
- non-crossing matchings
- non-crossing perfect matchings
- ...
- *[your favorite straight-line geometric graph structure]*

can it have?

We first consider the more popular variants - those with new works studying them every several years.

GRAPH TYPE	LOWER BOUND	REFERENCE	UPPER BOUND	REFERENCE
Plane Graphs	$\Omega(41.18^N)$	[AHHHKV]	$O(187.53^N)$	[SS12]
Triangulations	$\Omega(8.65^N)$	[DSST11]	30^N	[SS11]
Spanning Cycles	$\Omega(4.64^N)$	[GNT00]	$O(54.55^N)$	[SSW13]
Perfect Matchings	$\Omega(3.09^N)$	[AR15]	$O(10.05^N)$	[SW06]
Spanning Trees	$\Omega(12.52^N)$	[HM13]	$O(141.07^N)$	[HSSTW11; SS11]
Cycle-Free Graphs	$\Omega(13.61^N)$	[HM13]	$O(160.55^N)$	[HSSTW11; SS11]

Some less common variants:

Polynomial Method Lecture Notes #2

Polynomial Method Lecture Notes

Recent Comments

Incidences: L
Bo... on
Incidences: L
Bounds (par

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adamsheffer
The Two
Formulation
the Sz...



Brendan Mu
on The Two
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Min #Triangulations: $\Omega(2.43^n)$ $O(3.455^n)$

Perfect Matchings	$\Omega(3.09^N)$	[AR15]	$O(10.05^N)$	[SW06]
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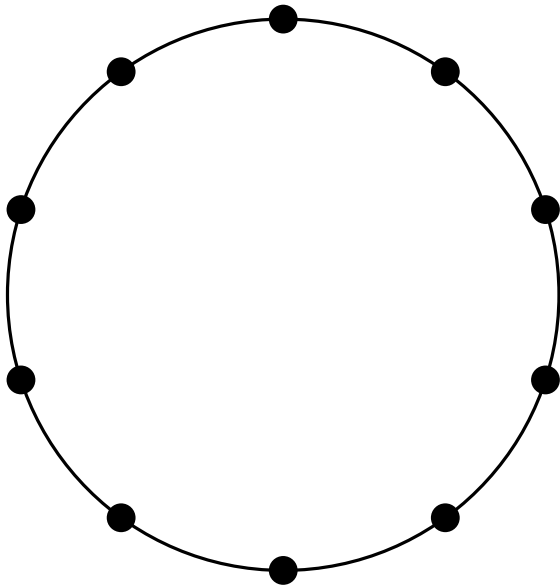
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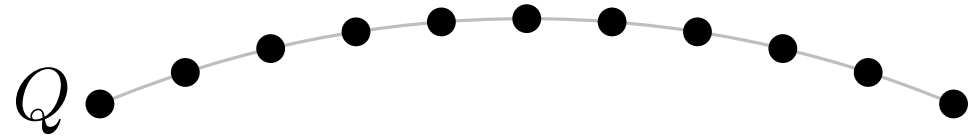
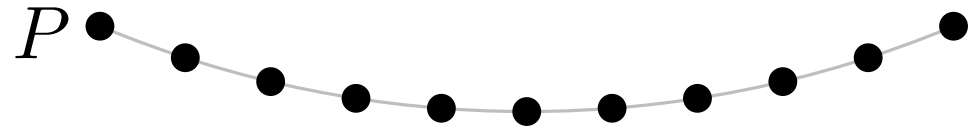
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convex position



double-chain

smallest possible number of perfect matchings: $\Theta^*(2^n)$

previous record: $\Theta^*(3^n)$

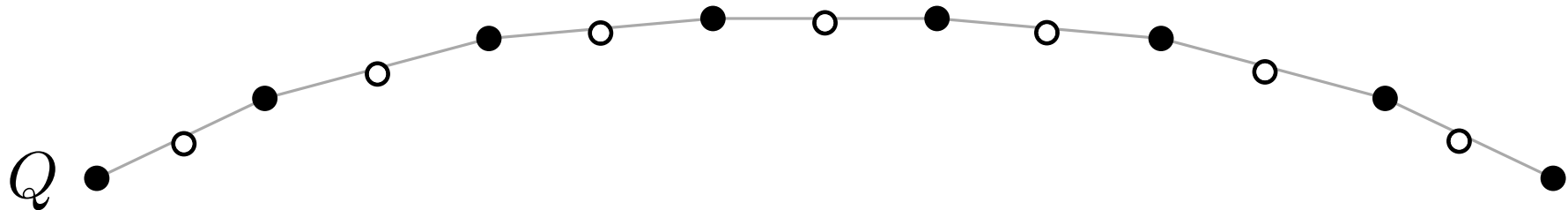
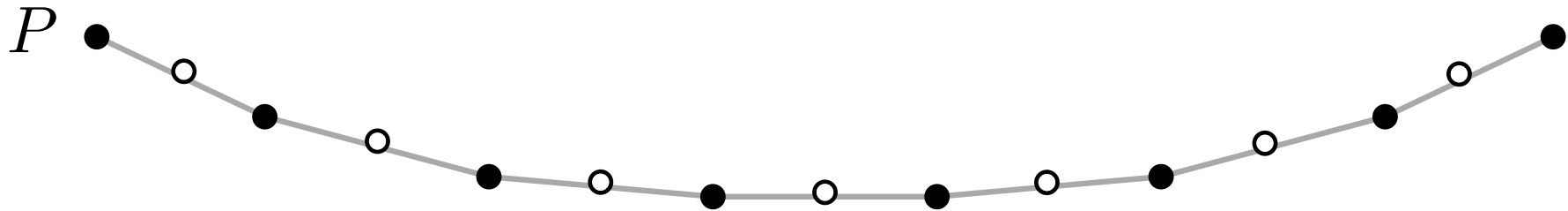
[García, Noy, Tejel 2000]

Upper bound: $O^*(10.06^n)$

[Sharir, Welzl 2006]

* = up to a polynomial factor

The Double-Zigzag Chain



$$C = \frac{2(1+x+x^3) - \sqrt{2(1+x+x^3)(1-2x-8x^2-3x^3+4x^4)}}{4x(1+x)(1+x+x^2)}$$

smallest singularity: $1 - 9x - 3x^2 = 0$

$$x_0 = \frac{\sqrt{93}}{6} - \frac{3}{2}$$

$$1/\sqrt{x_0} = \sqrt{6/(\sqrt{93} - 9)} \approx 3.0532$$

$$\#(\text{perfect matchings in } P \cup Q) = \Theta^*(3.0532^n)$$

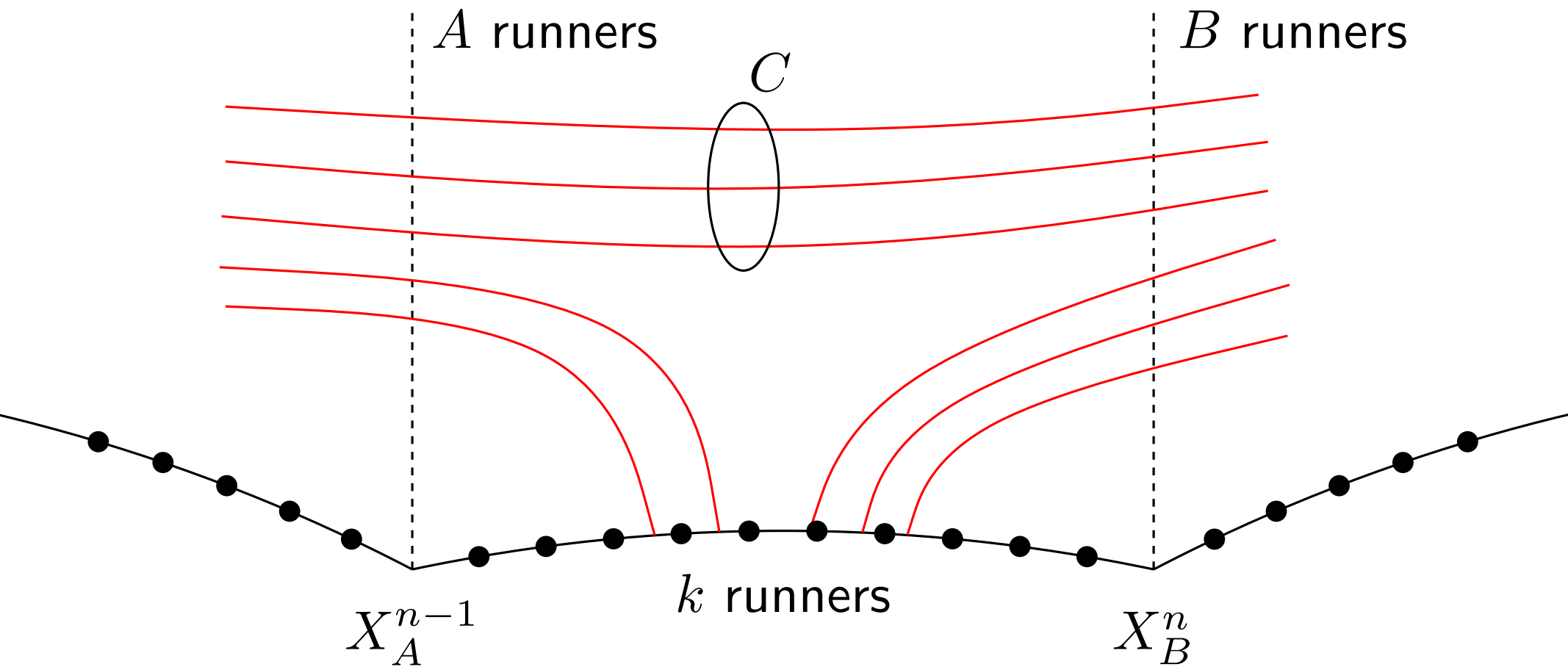
Longer Arcs

$$|P| = nr + 1$$



$$r = 8: \Theta^*(3.0930^n)$$

[joint work with Andrei Asinowski]



$X_B^n = \#$ possibilities after n arcs with B crossing runners

Example: $r = 5$

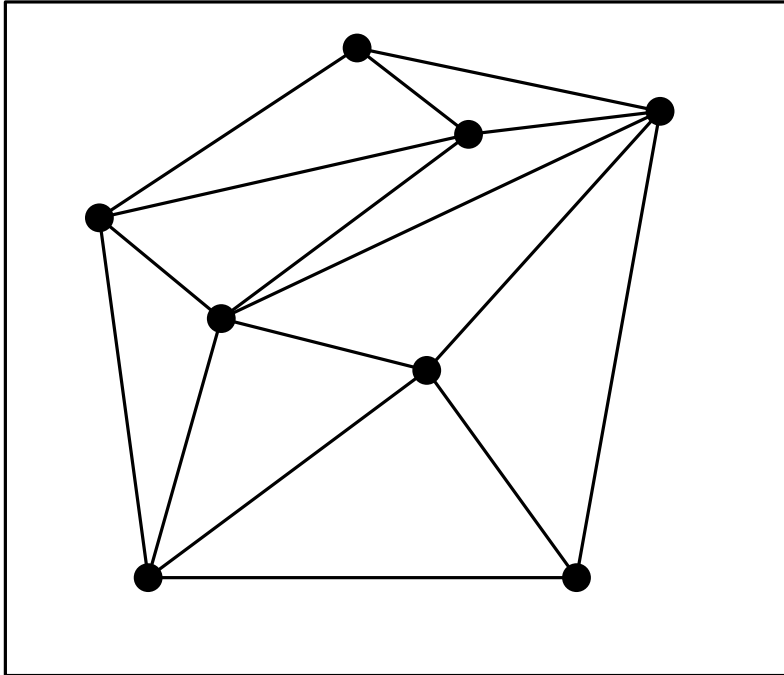
matrix for transforming $(X_0^{n-1}, X_1^{n-1}, X_2^{n-1}, \dots)$ into $(X_0^n, X_1^n, X_2^n, \dots)$

$$\begin{pmatrix} 10 & 30 & 30 & 20 & 5 & 1 & 0 & 0 & 0 & 0 & 0 & \dots \\ 30 & 40 & 50 & 35 & 21 & 5 & 1 & 0 & 0 & 0 & 0 & \dots \\ 30 & 50 & 45 & 51 & 35 & 21 & 5 & 1 & 0 & 0 & 0 & \dots \\ 20 & 35 & 51 & 45 & 51 & 35 & 21 & 5 & 1 & 0 & 0 & \dots \\ 5 & 21 & 35 & 51 & 45 & 51 & 35 & 21 & 5 & 1 & 0 & \dots \\ 1 & 5 & 21 & 35 & 51 & 45 & 51 & 35 & 21 & 5 & 1 & \dots \\ 0 & 1 & 5 & 21 & 35 & 51 & 45 & 51 & 35 & 21 & 5 & \dots \\ 0 & 0 & 1 & 5 & 21 & 35 & 51 & 45 & 51 & 35 & 21 & \dots \\ 0 & 0 & 0 & 1 & 5 & 21 & 35 & 51 & 45 & 51 & 35 & \dots \\ 0 & 0 & 0 & 0 & 1 & 5 & 21 & 35 & 51 & 45 & 51 & \dots \\ 0 & 0 & 0 & 0 & 0 & 1 & 5 & 21 & 35 & 51 & 45 & \dots \\ \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$

row sum = 271 \implies vectors grow like $271^n / \text{poly}(n)$

Counting, sampling, enumerating [V. Alvarez, R. Seidel 2013]

triangulation

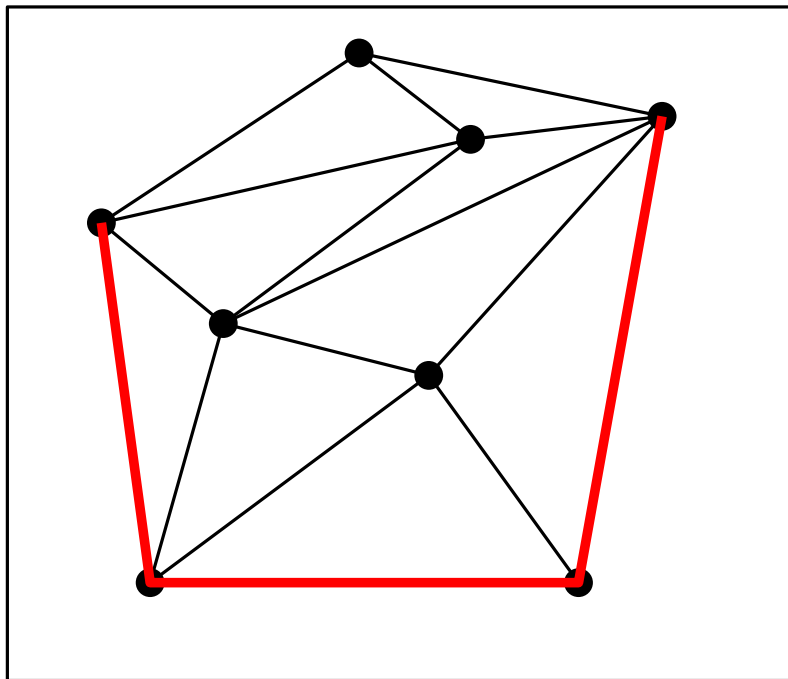


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triangulation



sequence of x -monotone paths

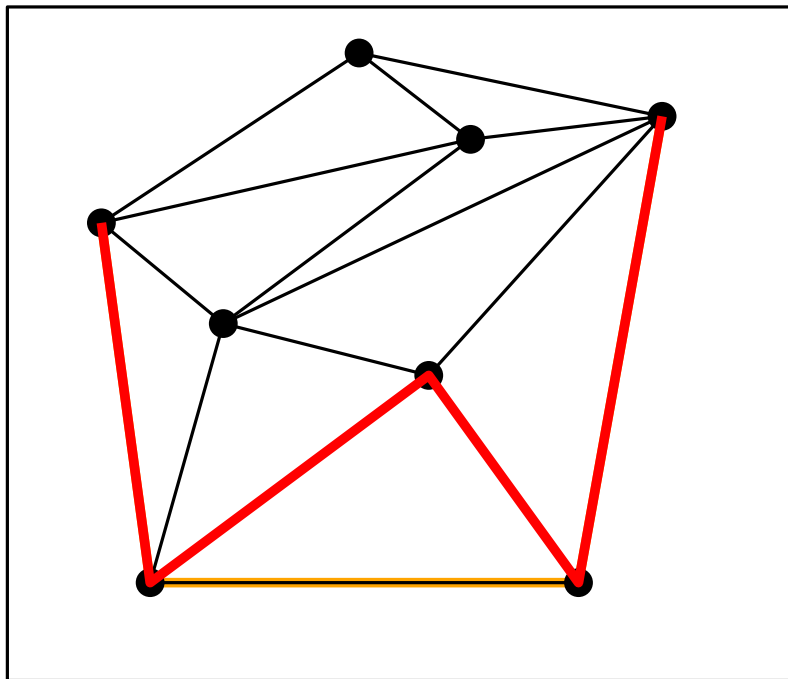


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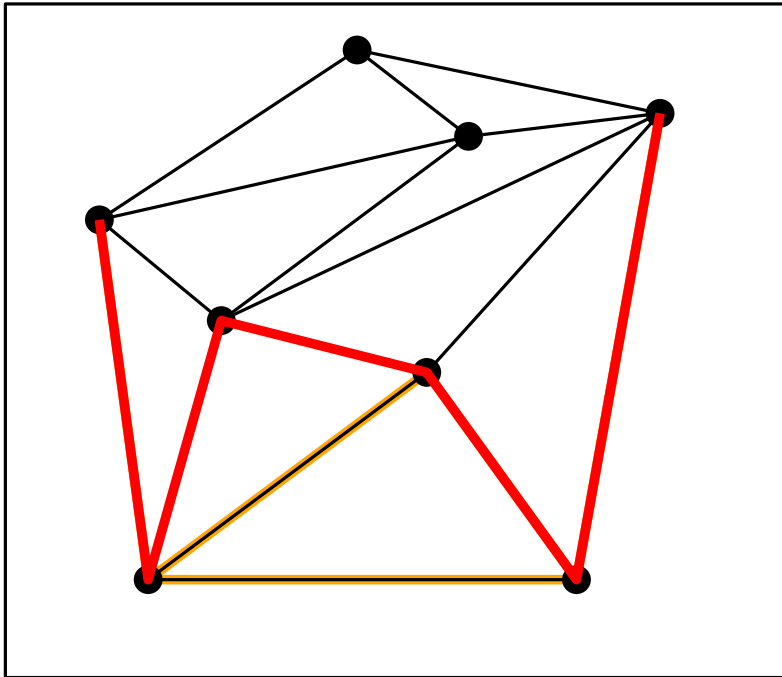


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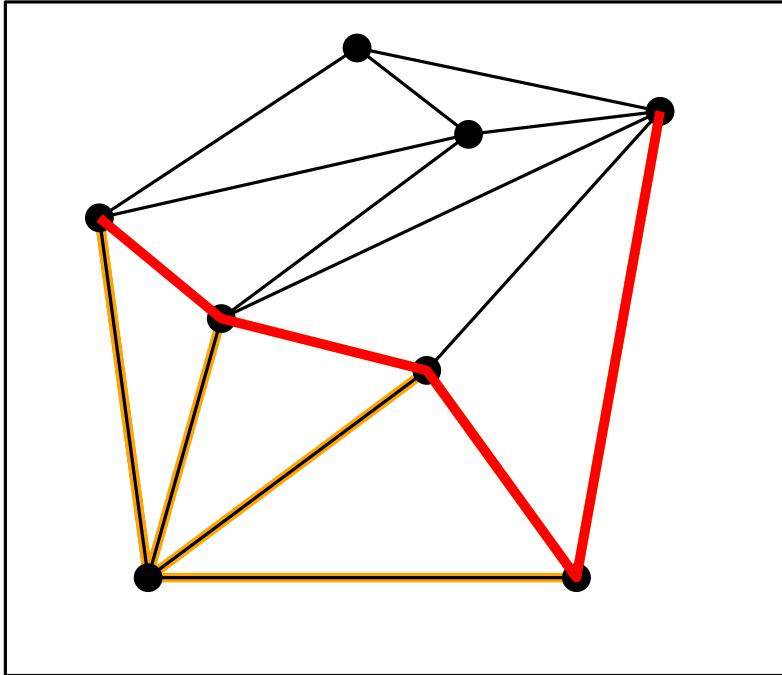


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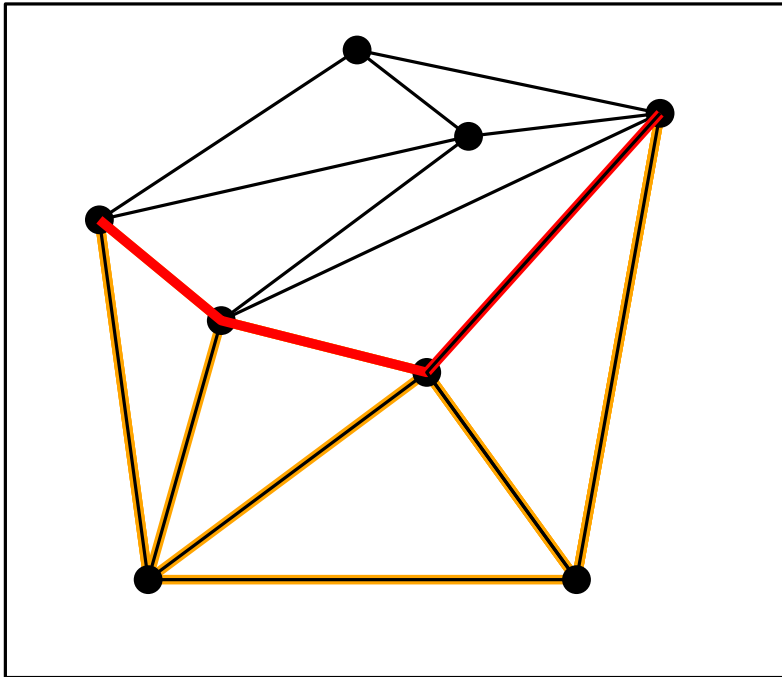


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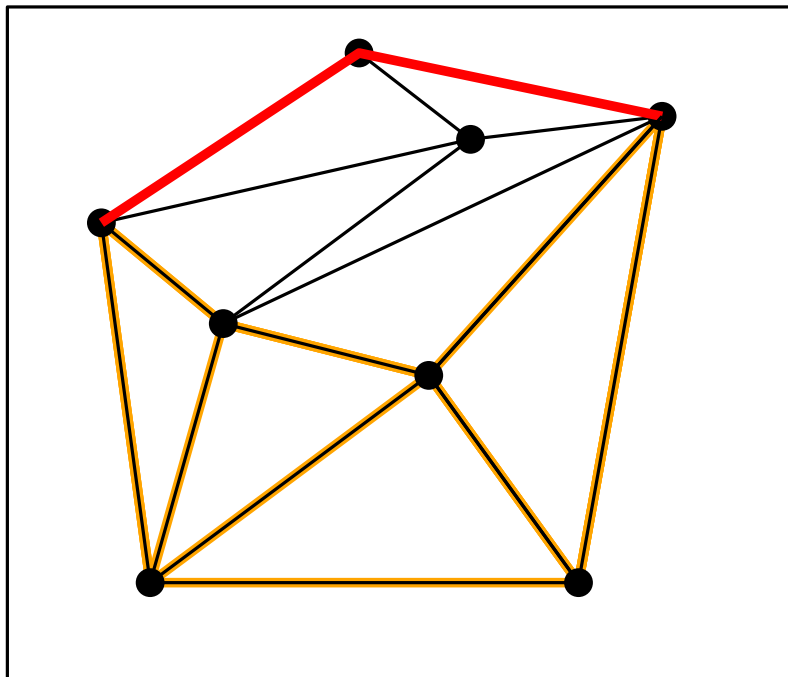


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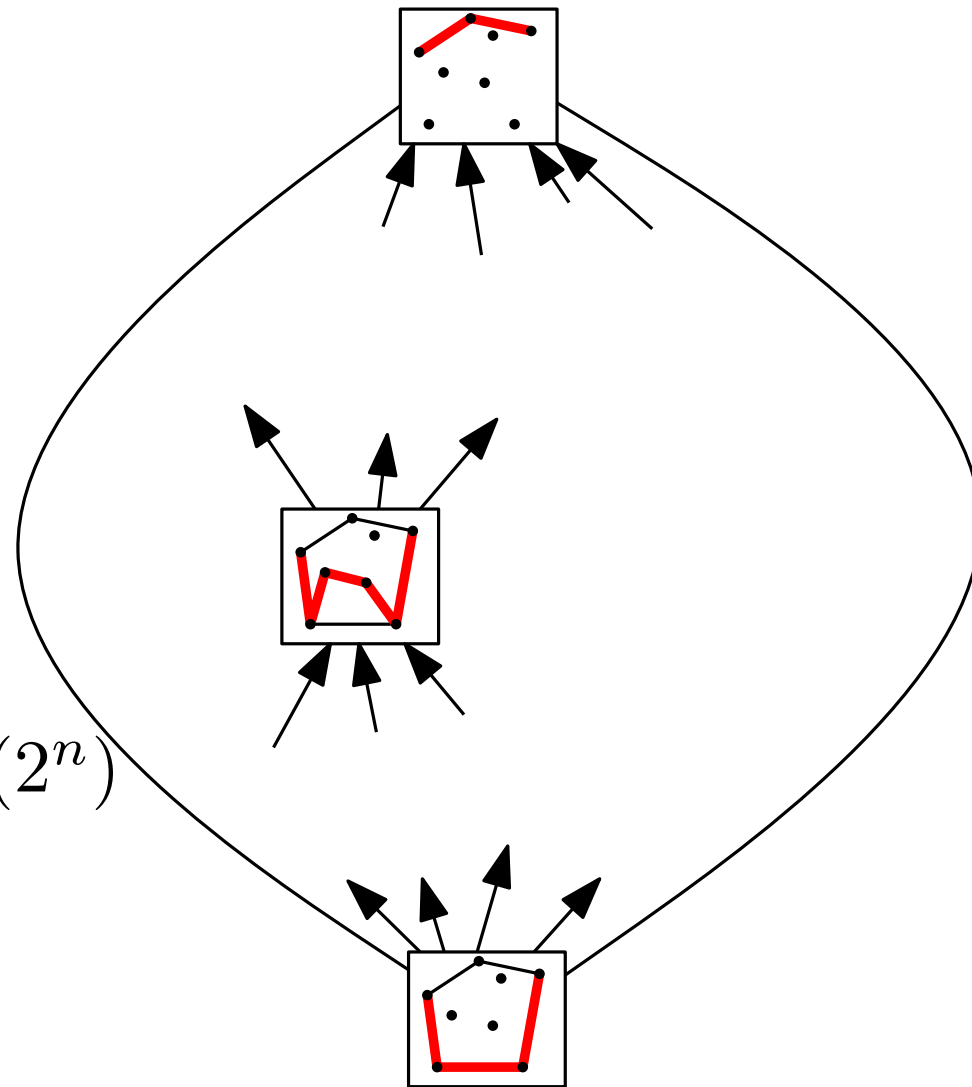
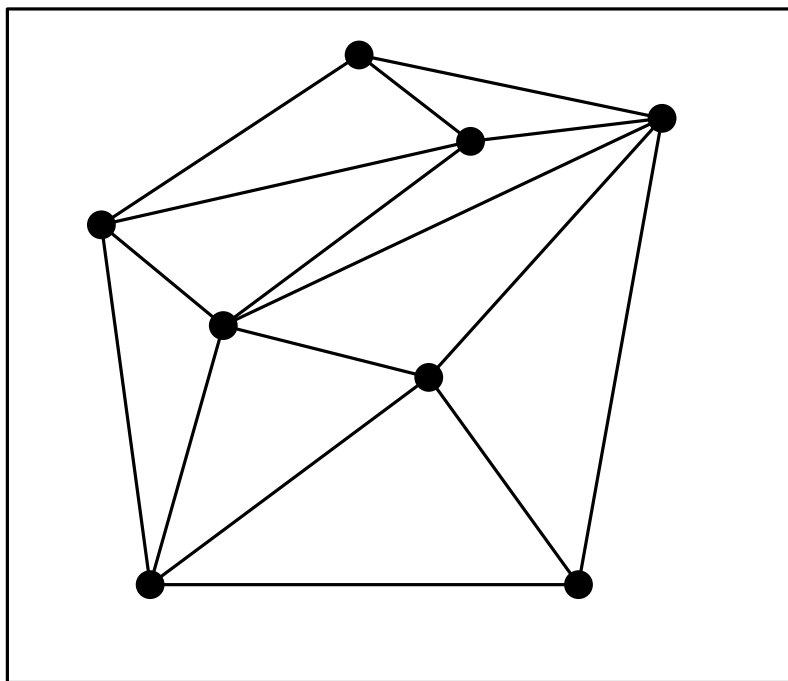


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triangulation

→

sequence of x -monotone paths



→ path in a DAG of size $O^*(2^n)$

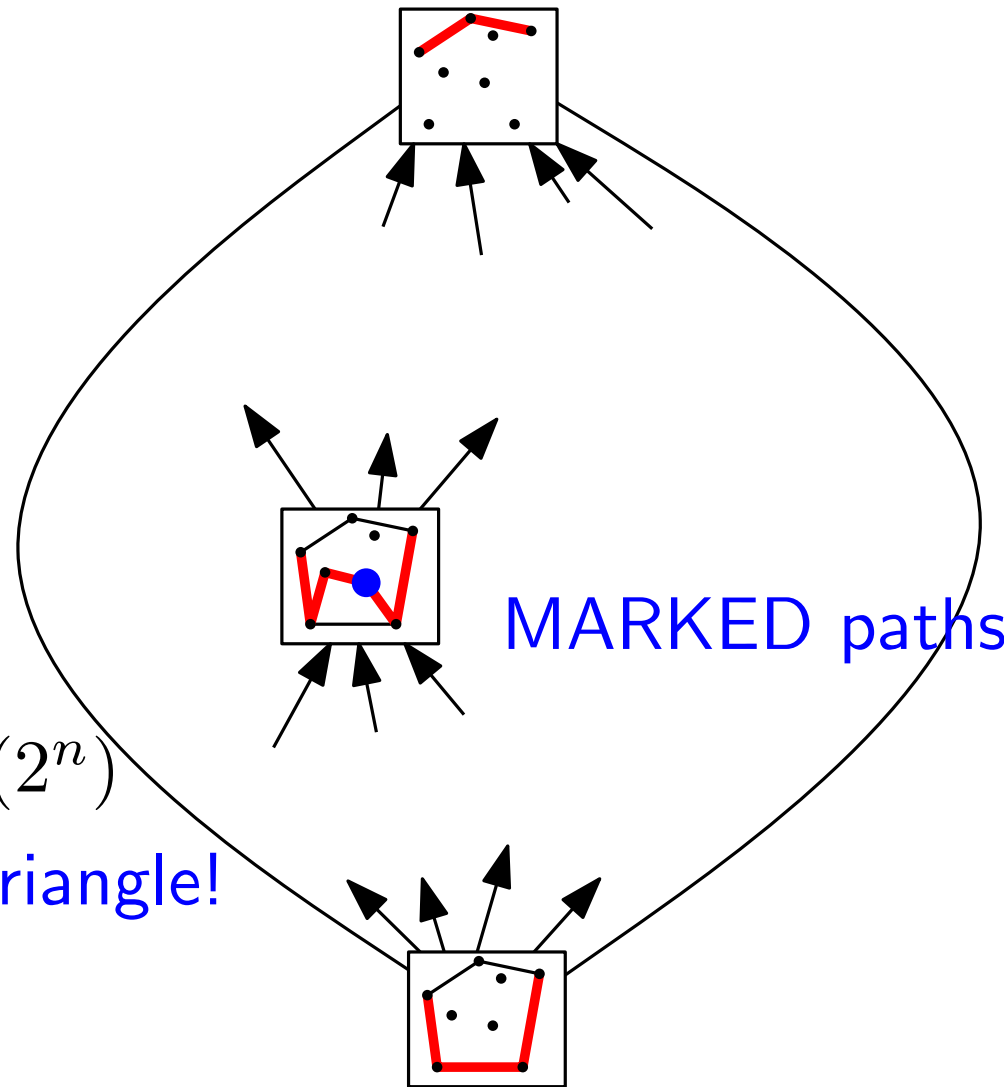
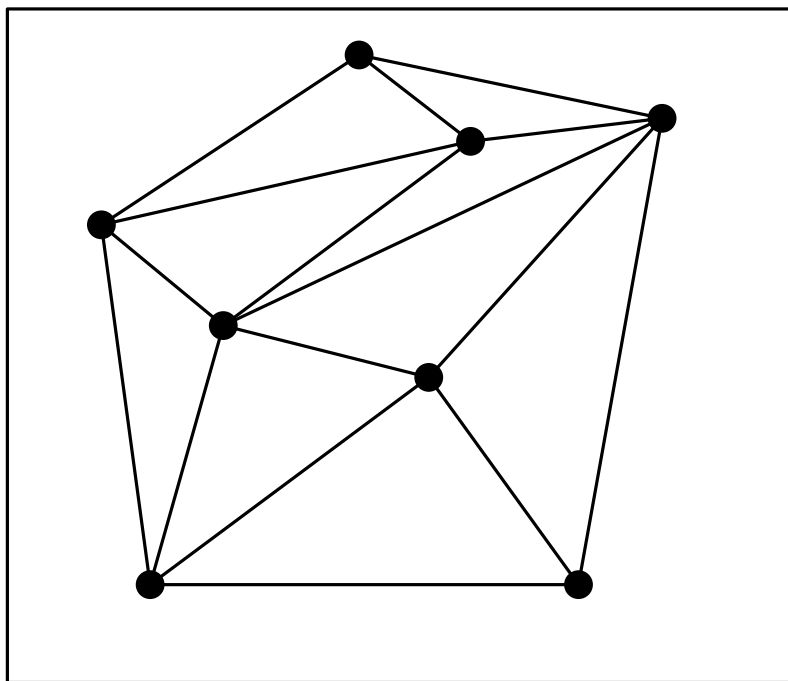
Counting Triangulations

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always choose the LEFTmost triangle!

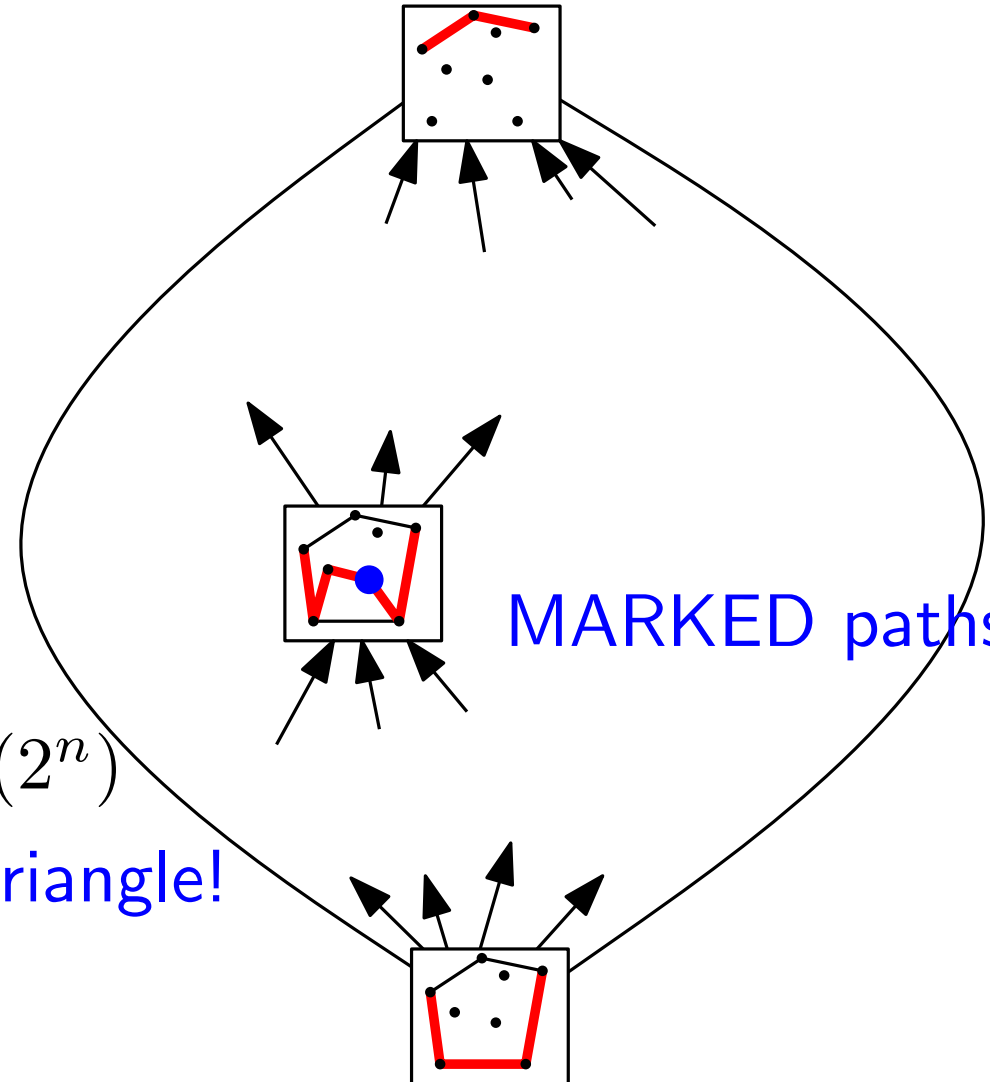
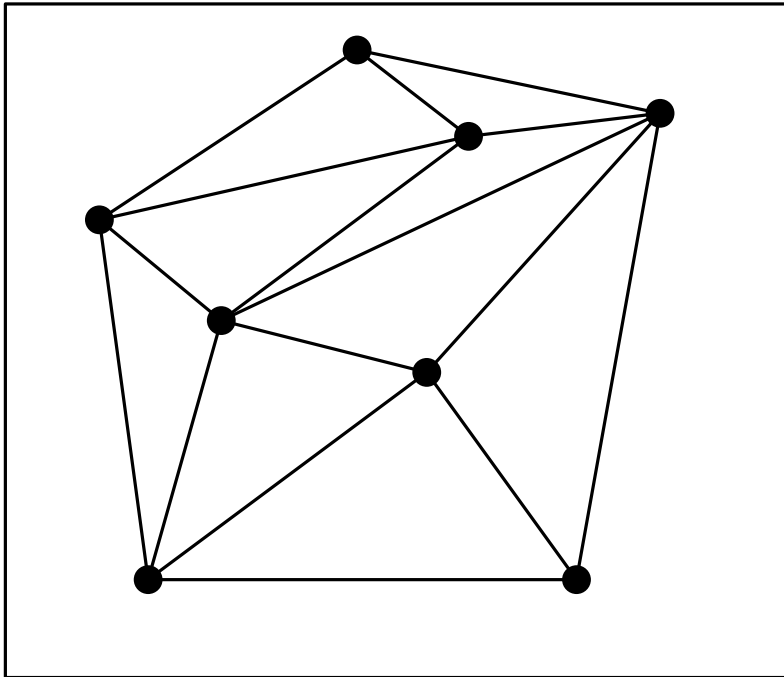
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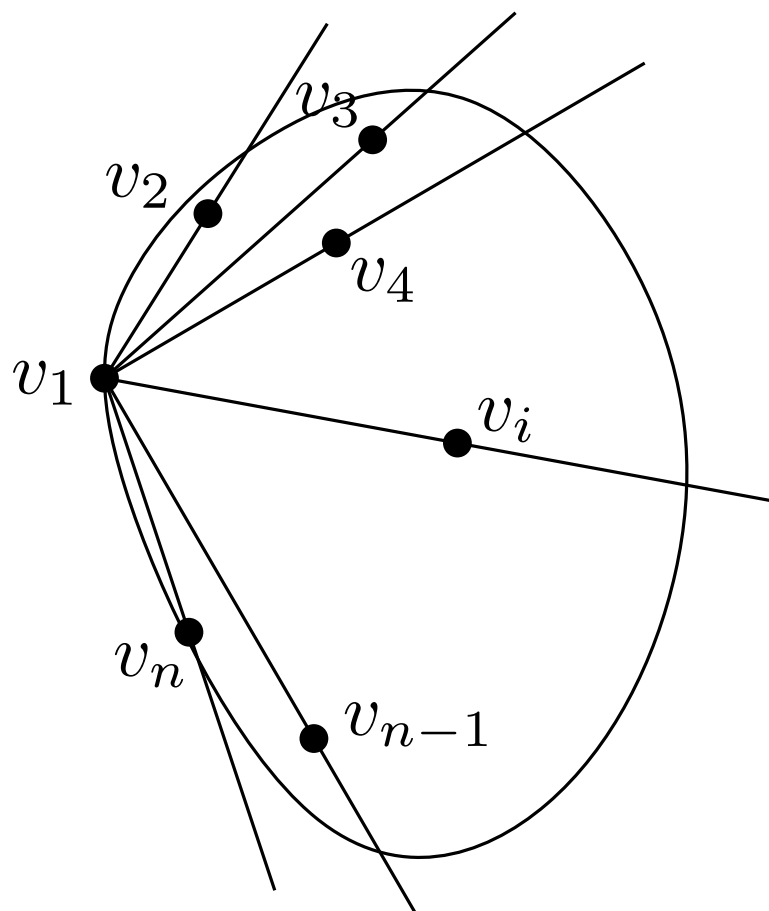
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$O(1)$ -delay enumeration,
with $O^*(2^n)$ preprocessing

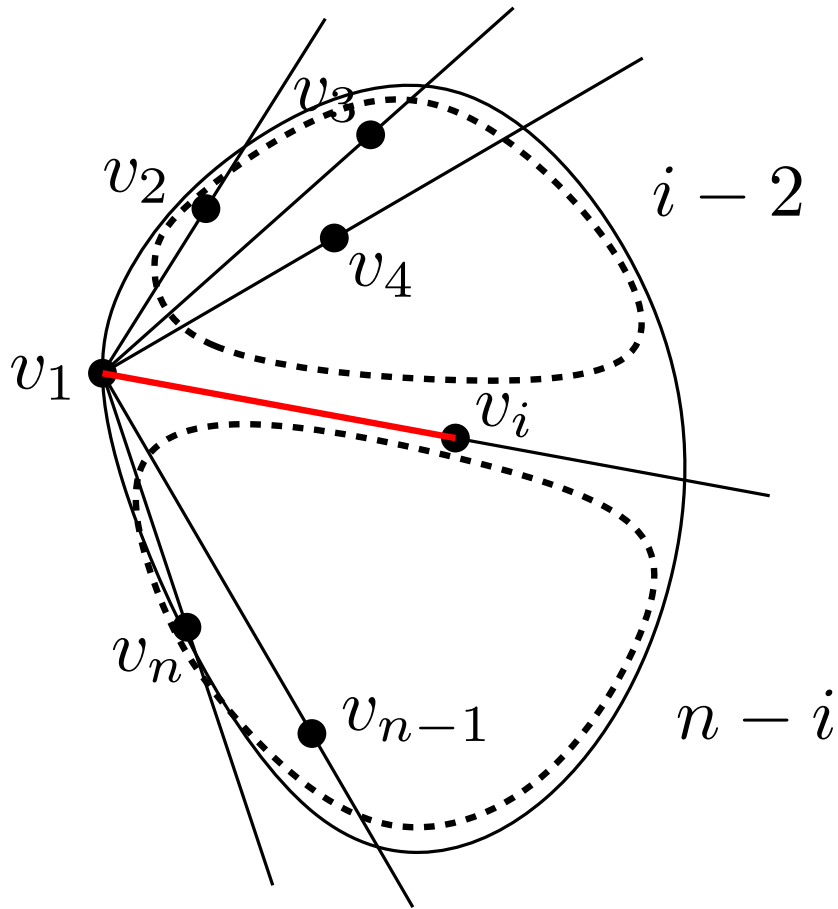
Extension to Perfect Matchings

Every point set has at least $\text{Catalan}(n/2) \sim 2^n$ perfect non-crossing matchings. [Manuel Wettstein 2014]



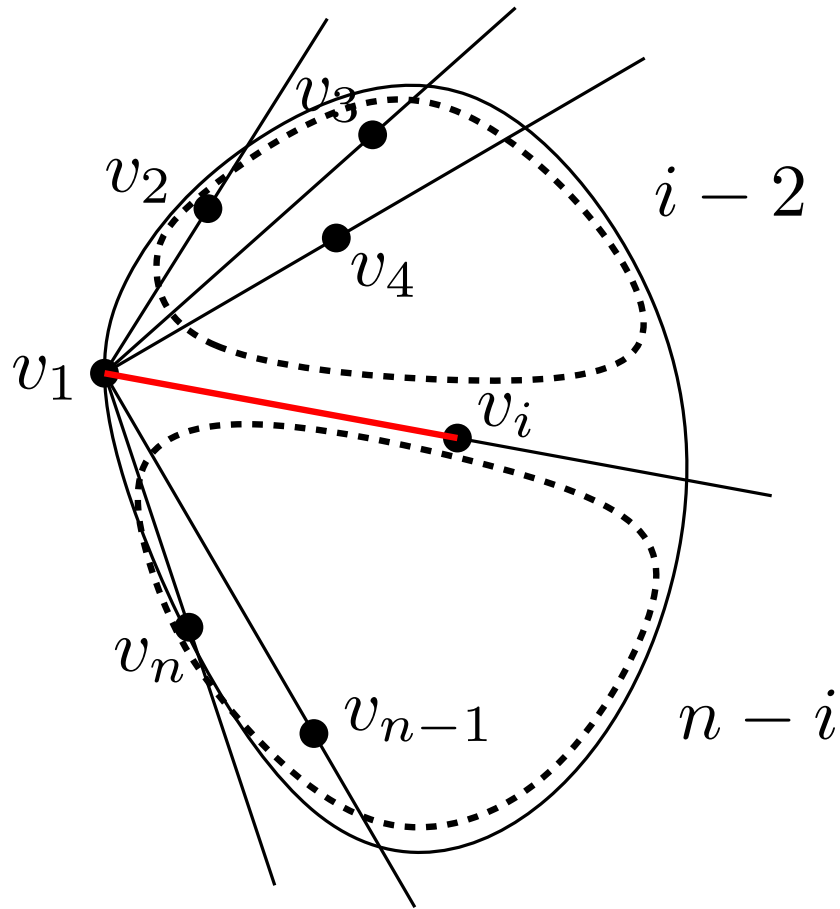
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(tight (almost only) for point sets in convex position)

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TRICK to achieve polynomial delay:

Output those “trivial” matchings while preparing the DAG.

(tight (almost only) for point sets in convex position)